HP StorageWorks Continuous Access EVA user guide performance estimator



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Continuous Access EVA performance estimator user guide

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About this guide

This guide provides information about:

- · Manually estimating replication performance for one outstanding write
- Determining the impact of distance on replication performance
- Using the Performance Estimator spreadsheet to estimate replication performance for one outstanding write
- Using the output from the manual calculation and the spreadsheet to determine the minimum expected
 performance in I/O operations per second (IOPS) for a particular I/O size, separation distance, and
 intersite link technology
- Estimating the ideal maximum throughput requirements for multiple outstanding writes

Intended audience

This guide is intended for customers who:

- Purchased the HP StorageWorks Continuous Access Enterprise Virtual Array (EVA)
- Want to estimate the effects of distance on applications that use HP StorageWorks Continuous Access EVA

Prerequisites

Prerequisites for using this product include knowledge of:

- SAN fabric configurations
- Disaster planning
- HP StorageWorks Enterprise Virtual Array (EVA)

Related documentation

In addition to this guide, pleas see other documents for this product:

- HP StorageWorks Continuous Access EVA administrator guide
- HP StorageWorks Continuous Access EVA planning guide
- HP StorageWorks SAN design reference guide

These and other HP documents can be found on an HP web site: http://www.docs.hp.com.

Document conventions and symbols

Table 1 Document conventions

Convention	Element
Medium blue text: Figure 1	Cross-reference links and e-mail addresses
Medium blue, underlined text (http://www.hp.com)	Web site addresses
Bold font	 Key names Text typed into a GUI element, such as into a box GUI elements that are clicked or selected, such as menu and list
Italics font	items, buttons, and check boxes Text emphasis
Monospace font	 File and directory names System output Code Text typed at the command-line
Monospace, italic font	Code variablesCommand-line variables
Monospace, bold font	Emphasis of file and directory names, system output, code, and text typed at the command-line



WARNING! Indicates that failure to follow directions could result in bodily harm or death.



CAUTION: Indicates that failure to follow directions could result in damage to equipment or data.



IMPORTANT: Provides clarifying information or specific instructions.



NOTE: Provides additional information.



TIP: Provides helpful hints and shortcuts.

HP technical support

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Collect the following information before calling:

- Technical support registration number (if applicable)
- Product serial numbers
- Product model names and numbers
- Applicable error messages
- Operating system type and revision level
- Detailed, specific questions

For continuous quality improvement, calls may be recorded or monitored.

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HP-authorized reseller

For the name of your nearest HP-authorized reseller:

- In the United States, call 1-800-345-1518.
- Elsewhere, visit http://www.hp.com and click Contact HP to find locations and telephone numbers.

Helpful web sites

For third-party product information, see the following vendor web sites:

- http://www.hp.com
- http://www.hp.com/go/storage
- http://www.hp.com/support/
- http://www.docs.hp.com

1 Introduction

When you are designing an HP StorageWorks Continuous Access solution, it is important to understand and plan the distance between the two sites to ensure optimum performance and disaster tolerance. A shorter distance between sites will improve performance, but it will also increase the risk of a disaster affecting both sites (reduced disaster tolerance). A longer distance between sites will reduce the risk of disaster affecting both sites (improved disaster tolerance), but it will also decrease performance. The key to a successful solution is creating a balance between performance and protection.

This chapter describes the replication process and the effect of bandwidth on replication performance. Using these two components, you can use this guide to estimate:

- The performance capabilities of an intersite link technology.
- The link requirements for a given level of performance.

Replication process

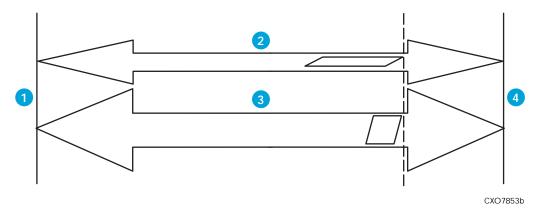
To complete one replication I/O, Continuous Access EVA requires one round trip through the intersite link as follows:

- 1. The source array sends the data to the destination array.
- 2. The destination array receives and stores the data in its local cache and returns an acknowledgement to the source array.

Bandwidth and replication performance

HP Continuous Access products can move data at extreme distances. However, the inherent propagation delays in data transfers can affect the amount of time it takes to complete replication I/O. At these extreme distances, bandwidth is not typically the limiting factor in replication performance; rather it is the delay imposed by that distance. At short distances, bandwidth affects performance more than the distance-induced delay does.

To better understand the relationship between bandwidth, distance, and performance, consider a low bandwidth intersite link and a high bandwidth intersite link that are both moving I/O packets containing the same amount of data (Figure 1). The packets are moving from site A to site B (left to right), with the leading edge of both packets arriving at site B at the same time. The difference in link size (bandwidth) enables the high bandwidth link to complete delivery of the packet (loading and unloading) before the low bandwidth link.



- 1 Site A
- 2 T3 link (44.5 Mb/s)
- 3 OC3 link (155 Mb/s)
- 4 Site B

Figure 1 Distance and I/O rate

<u>-</u>

TIP: For a visual reference, a nanosecond of fiber is approximately eight inches long. A bit is four inches long at 2 Gb/s and eight inches long at 1 Gb/s. Therefore, a nanosecond of fiber contains two bits at 2 Gb/s and one bit at 1 Gb/s. At 100 Mb/s, a bit is 80 inches and at 10 Mb/s, a single bit is 800 inches.

2 Manually estimating performance

This chapter describes how to manually estimate the time needed to complete one I/O replication across a directly connected (or zero distance) intersite link. It also describes how to add the impact of distance to the calculation.

Topics include:

Using the formula without distance, page 11

Adding distance, page 12

Calculating performance for one outstanding write, page 12

Using the formula without distance

To calculate performance based on link bandwidth, use the following formula:

$$Y = mX + b$$

Where

- y is the time (in milliseconds) to complete one outstanding replication I/O.
- m is the slope of the line representing the link bandwidth.
- x is the size of the transfer in kilobytes.
- mx is the is the additional time required to complete a transfer (or data packet) larger than 512 bytes for a given link bandwidth.
- b is the amount of time required to replicate a single 512 byte write (replication and conversion overhead). This is the y-intercept of the line representing the link bandwidth when x is zero.

Use the values in Table 2 for the appropriate intersite link technology to complete this formula.

NOTE: The values in Table 2 were derived experimentally.

Table 2 Intersite link values

Technology	Bandwidth (Mb/s)	Intercept (b)	Slope (m)
2 Gb/s Fibre Channel	2000	0.3416	0.0268
1 Gb/s Fibre Channel	1000	0.3991	0.0332
1 GbE (Gigabit Ethernet) IP	1000	0.4130	0.0338
OC3 IP	155.5	0.3901	0.0758
E4 IP	139.3	0.3876	0.0818
100 Mb/s IP	100	0.3802	0.1052

Table 2 Intersite link values (continued)

Technology	Bandwidth (Mb/s)	Intercept (b)	Slope (m)
T3 IP	44	0.3530	0.2070
E3 IP	34.304	0.3340	0.2666
10 Mb/s IP	10	0.2893	0.8872
E1 IP	2.048	1.1748	4.4434
T1 IP	1.54	1.5557	5.9422

Adding distance

To obtain a more realistic estimate of performance, add the effects of distance, which is measured by the time required for the leading edge of the data to traverse the link. Estimate this time using a speed of 2×10^8 meters per second for light in standard fiber-optic cable, which is equal to 5 microseconds per kilometer. For example, if the intersite link distance is 10,000 km (6,200 mi), the one-way time interval is 50,000 microseconds, or 0.05 seconds. This is approximately 25 times the average rotational latency of a 15,000 rpm disk drive at 0.002 seconds.

Therefore, the actual time to move data from the source array to the destination array is the round-trip distance, plus the time to load and unload the data for the link size. These trips consist of one small control packet and one large data packet and add 10 microseconds of latency (per kilometer) to the one-way intersite link distance to complete each replication write. Based on the distance between the two sites, this latency is added to the previously calculated time to complete a zero distance replication write.

Calculating performance for one outstanding write

During synchronous replication, an application performs one write and waits for that write to complete before issuing the next outstanding write. This section describes how to calculate performance for one outstanding write that contains a 32 KB data packet. Based on the one-way network latency of 1 ms, the local and remote sites are 200 km (125 mi) apart.

Topics include:

- Calculating replication time, page 13
- Calculating time to load, page 13
- Calculating time to complete, page 14
- Determining maximum I/Os per second, page 15

Calculating replication time

To measure the performance of this outstanding write:

- 1. Use the values in Table 2 to calculate the loading time for a 32 KB data packet. Multiply the slope by the size of the transfer m times x).
- Calculate and add the time it takes for a round trip. If the intersite link distance is 200 km, the transfer (roundtrip) latency is calculated as follows (applies to all intersite link types):

```
200 km x 5 ms/km x 2 = 2 ms
```

3. Add the replication overhead from Table 2 (the Intercept column).

The sum is the total time it takes to replicate a 32 KB data packet for a single outstanding write.

Calculating time to load

The *time-to-load* value is defined as the length of the data packet in time (the time to load and unload it). Using the intersite link bandwidth and slope values from Table 2, Table 3 shows the time-to-load value (step 1) for a 32 KB data packet using different intersite link technologies.(

Table 3 Time-to-load value

Intersite link technology	Slope	Time to load (ms)
2 Gb/s Fibre Channel	0.0268	0.86
1 Gb/s Fibre Channel	0.0332	1.06
1 GbE (Gigabit Ethernet) IP	0.0338	1.08
OC3 IP	0.0758	2.43
E4 IP	0.0818	2.62
100 Mb/s IP	0.1052	3.37
T3 IP	0.2070	6.62
E3 IP	0.2666	8.53
10 Mb/s IP	0.8872	28.39
E1 IP	4.4434	142.19
T1 IP	5.9422	190.15

Calculating time to complete

The *time-to-complete* value is defined as the time to complete one I/O replication across a given technology for a given packet size. It is the time for a host to send a synchronous replication write and receive an acknowledgement of completion. Table 4 show the *time-to-complete* value for a 32 KB data packet (the sum of step 2 and step 3) by adding transfer latency and replication overhead to the results in Table 3.

Table 4 Time-to-complete value

Intersite link technology	Time to load (ms)	+ Transfer latency (ms)	+ Overhead (ms)	= Time to complete I/O (ms)
2 Gb/s Fibre Channel	0.86	2	0.3416	3.20
1 Gb/s Fibre Channel	1.06	2	0.3991	3.46
1 GbE (Gigabit Ethernet) IP	1.08	2	0.4130	3.49
OC3 IP	2.43	2	0.3901	4.82
E4 IP	2.62	2	0.3876	5.01
100 Mb/s IP	3.37	2	0.3802	5.75
T3 IP	6.62	2	0.3530	8.98
E3 IP	8.53	2	0.3340	10.87
10 Mb/s IP	28.39	2	0.2893	30.68
E1 IP	142.19	2	1.1748	145.36
T1 IP	190.15	2	1.5557	193.71

Determining maximum I/Os per second

You can use the time-to-complete value to determine the maximum number of synchronous replication writes that can be completed every second, assuming the next I/O starts immediately after the current I/O completes. To do so, you must invert the time to complete value into I/Os per second (IOPS). Table 5 shows the results of the inversion for a 32 KB data packet, which were calculated as follows:

- Throughput is the packet size multiplied by the I/O rate (IOPS).
- Bandwidth used is the peak bandwidth of the link, divided by the throughput and then multiplied by 100.

NOTE: Because most applications produce multiple asynchronous I/Os, Table 5 shows the minimum expected performance.

Table 5 Bandwidth used for a 32 KB write

Intersite link technology	Approximate IOPS	Throughput (Mb/s)	Approximate single stream bandwidth used
2 Gb/s Fibre Channel	312.58	100.03	5.0%
1 Gb/s Fibre Channel	288.89	92.45	9.2%
1 GbE (Gigabit Ethernet) IP	286.16	91.57	9.2%
OC3 IP	207.65	66.45	42.9%
E4 IP	199.79	63.93	46.0%
100 Mb/s IP	174.02	55.69	55.7%
T3 IP	111.40	35.65	79.2%
E3 IP	92.04	29.45	86.6%
10 Mb/s IP	32.59	10.43	100.0%
E1 IP	6.88	2.20	100.0%
T1 IP	5.16	1.65	100.0%

Using the Performance Estimator

This chapter describes how to use the Performance Estimator, a Microsoft Excel-based spreadsheet, to perform the calculations described in "Manually estimating performance" on page 11.

Topics include:

- Accessing the spreadsheet, page 17
- Entering one-way latency, page 18
- Entering data packet size, page 19
- Reading the results, page 19
- Comparing calculation results, page 21

Accessing the spreadsheet

The Performance Estimator spreadsheet is available from the following web site:

http://h18006.www1.hp.com/products/storage/software/conaccesseva/index.html

Select Related information and then select HP StorageWorks Continuous Access EVA Performance **Estimator** to download the spreadsheet.

Figure 2 shows the default view of the spreadsheet when you open it. When you enter values in the latency and size fields, the spreadsheet calculates the bandwidth for:

- **2 Gb/s fiber**—Direct fiber or wave dimension multiplexing (WDM) only.
- 1 Gb/s fiber—Direct fiber over WDM only.
- FC-SONET or FC-IP—See Table 2 on page 11 for a complete list of the intersite link technologies in this category.

The 2 Gb/s and 1 Gb/s intersite links include long-distance direct fiber connections using longwave or very long distance GBIC/SFP, of either coarse- or dense-wave division multiplexing (CWDM or DWDM).



TIP: Place the mouse over any red triangle in the spreadsheet to view additional information for that field.

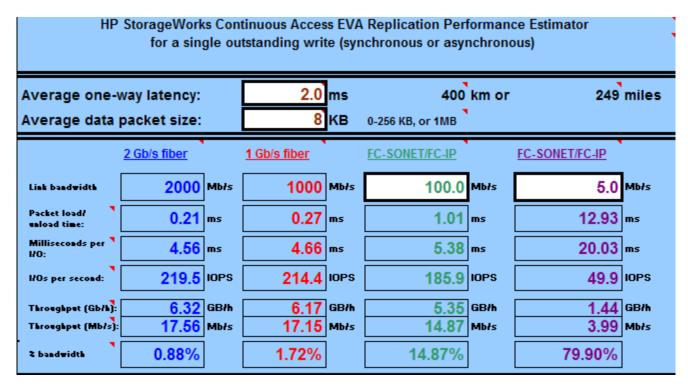


Figure 2 Performance Estimator spreadsheet

Entering one-way latency

To begin using the spreadsheet, enter the appropriate value in the Average one-way latency field. This is the distance between the source and destination arrays in milliseconds. Use the ping command to determine round-trip latency and then divide the result in half. You can also estimate latency using one of the following calculations:

- For point-to-point networks, multiply the driving distance between sites by 1.5.
- For routed networks, multiply the driving distance between sites by 2.25.

When you enter a value and either press **Enter** or click outside the cell, the latency is translated into kilometers and miles. Figure 3 shows three examples.

NOTE: The maximum value for this field is 100milliseconds.

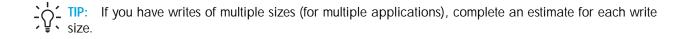
Average one-way latency:	4.0 ms	800 km or	497 miles
Average one-way latency:	10.0 ms	2,000 km or	1,243 miles
Average one-way latency:	15.0 ms	3,000 km or	1,864 miles

Figure 3 Latency examples

Entering data packet size

The other input is the average size of the application write being replicated to the destination array. The values for this field are 0–256 KB, or 1 MB for a full copy. You can estimate the size for the type of applications you are using. Some examples are:

- 4 KB (such as Microsoft Exchange)
- 8 or 16 KB (small databases, such as Microsoft SQL)
- 32 or 64 KB (large databases, such as Oracle)



Reading the results

Using the information you enter, the Replication Performance Estimator spreadsheet calculates how latency and write size affect the bandwidth of the applicable intersite link technology.

For all intersite link technologies listed (2 Gb/s fiber, 1 Gb/s fiber, and FC-SONET or FC-IP), the bandwidth results display as follows:

- Link bandwidth—The amount of bandwidth (in Mb/s)
- Packet load/unload time—The length of time (in milliseconds) to move the data packet on to and off of the intersite link. It is based on the size of the outstanding write and the link bandwidth.
- ms per I/O—The length of time to complete one synchronous I/O across the link. Starting with the
 intersite latency and packet transmit time formulas above. It is the summation of packet transmit time,
 plus twice the one-way latency, plus all conversion overhead.
- I/O per second—The inverse of milliseconds per I/O. The maximum number of synchronous I/O per second for a single data stream. For synchronous I/O, the next I/O does not start until the one in progress has finished. The maximum rate can be achieved only when the next write starts immediately after completion of the previous write. In a real-world environment in which writes are generated in a pseudo-random fashion, the expected peak is approximately 70% of the theoretical peak. The average rate typically does not exceed 50% of the theoretical peak. Table 5 on page 15 shows the I/Os per second of various SAN technologies.
- Throughput (GB/h)—The transfer rate based on one hour of I/Os per second, multiplied by the data packet size.
- Throughput (Mb/s)—The transfer rate based on one second of I/Os per second multiplied by the data packet size. Add these numbers for each I/O stream to get estimated bandwidth requirements.
- % bandwidth—The estimated bandwidth required for this single replication. It is based on the
 estimated throughput and the link bandwidth. Due to the mathematics of the model, the value may
 exceed 100%.

Figure 4 shows the results for latency of 5 milliseconds and data packet size of 32 KB.

	2 Gb/s fiber	•	1 Gb/s fiber	•	FC-SONET/FC-IP		FC-SONET/FC-IP	
Link bandwidth	2000	Mbłs	1000	Mbłs	100.0	Mbłs	5.0 мь	ols
Packet load/ unload time:	0.86	ms	1.06	ms	4.05	ms	51.73 ms	F
Milliseconds per 170:	11.20	ms	11.46	ms	14.42	ms	64.83 ms	F
I/Os per second:	89.3	IOPS	87.2	IOPS	69.3	IOPS	15.4 IOF	PS
Throughput (Gb/h)	10.29	GB/h	10.05	GB/h	7.99	GB/h	1.78 GB	3/h
Throughput (Mb/s)	28.57	Mb/s	27.92	Mb/s	22.19	Mbłs	4.94 мь	ols
2 bandwidth	1.43%		2.79%		22.19%		98.73%	

Figure 4 Sample results

Altering the link bandwidth

You cannot change the link bandwidth fields for 2 Gb/s and 1 Gb/s fiber. You can, however, alter the link bandwidth for both FC-SONET/FC-IP fields. This enables you to enter a value for the specific intersite link technologies you are using. The default values are 100 Mb/s (100 Mb/s Ethernet IP ISL) and 5 Mb/s (the minimum supported bandwidth for each ISL in a dual ISL connection). Figure 4 shows the results based on these default values.

If you want to see the results for the OC3 IP intersite link, enter 155.5 in the Link bandwidth field for the first FC-SONET/FC-IP column. Figure 5 shows the results.

NOTE: The comment for the FC-SONET/FC-IP column lists the value to enter for the corresponding ISL.

	2 Gb/s fiber		1 Gb/s fiber		FC-SONET/FC-IP		FC-SONET/FC-IP	
Link bandwidth	2000	Mbłs	1000	Mbłs	155.5	Mbłs	5.0 Mb/	ls
Packet load/ unload time:	0.86	ms	1.06	ms	2.79	ms	51.73 ms	
Milliseconds per 1/0:	11.20	ms	11.46	ms	13.16	ms	64.83 ms	
I/Os per second:	89.3	IOPS	87.2	IOPS	76.0	IOPS	15.4 IOP	S
Throughput (Gb/h):	10.29	GB/h	10.05	GB/h	8.75	GB/h	1.78 GB/	/h
Throughput (Mb/s)	28.57	Mb/s	27.92	Mb/s	24.31	Mbłs	4.94 мы	ls
2 bandwidth	1.43%		2.79%		15.63%		98.73%	

Figure 5 Sample results for OC3 IP intersite link

As another example, you want to see the bandwidth results for the T3 IP intersite link. Change the Link bandwidth value for the first FC-SONET/FC-IP column to the default (100) and enter 44 in the second FC-SONET/FC-IP column. Figure 6 shows the results.



NOTE: HP recommends that you begin with the default values for the FC-SONET/FC-IP fields. As your analysis proceeds, adjust one value to compare with the default value.

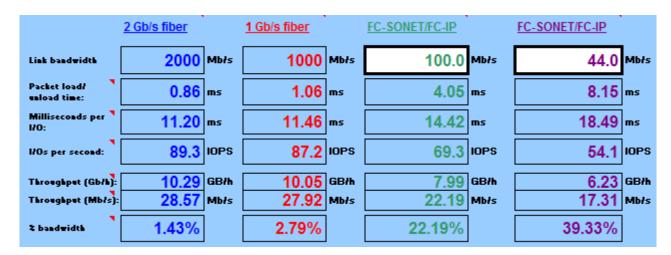


Figure 6 Sample results for T3 IP intersite link

Comparing calculation results

Figure 7 shows that latency of 1 millisecond and data packet size of 32 KB consumes approximately 50% of a 100 Mb/s FC-IP intersite link. The average load on any link must not exceed 40% of its rated capacity, and the peak loading must not exceed 45% of rated capacity. This limitation allows I/O from a failed link or fabric to run on the nonfailed fabric or link without causing additional failures by overloading the active fabric.

The result of the manual calculation for the same latency and data packet size is 55.77% (see Table 5 on page 15).



NOTE: There may be a difference between the results you generate manually and those you generate using the spreadsheet. The Performance Estimator spreadsheet uses a mathematical model based on the using the spreadsheet. The Performance Estimator spreadsheet uses a mathematical model based on the IP link bandwidth to estimate the slope and intercept of the line that is used in the first two rows of calculations.

Average one-way latency:			1.0	ms	200 km or 124		miles	
Average data	Average data packet size:			KB	0-256 KB, or 1MB			
	2 Gb/s fiber		1 Gb/s fiber		FC-SONET/FC-IP		FC-SONET/FC-IP	
Link bandwidth	2000	Mbłs	1000	Mb/s	100.0	Mbłs	44.0	Mbłs
Packet load/ unload time:	0.86	ms	1.06	ms	4.05	ms	8.15	ms
Milliseconds per 170:	3.20	ms	3.46	ms	6.42	ms	10.49	ms
I/Os per second:	312.6	IOPS	288.9	IOPS	155.8	IOPS	95.3	IOPS
Throughput (Gb/h):	36.01	GB/h	33.28	GB/h	17.94	GB/h	10.98	GB/h
Throughput (Mb/s):	100.03	Mbłs	92.45	Mbłs	49.84	Mbłs	30.51	Mbłs
2 bandwidth	5.00%		9.24%		49.84%		69.33%	

Figure 7 Results for a 32 KB write

4 Determining maximum bandwidth

The previous chapters showed results for one outstanding write. However, most applications do not issue one write and wait for it to complete before immediately sending the next one. Instead, multiple writes are replicated simultaneously (asynchronous replication), which requires that you determine the effect of multiple writes from a single application. This chapter describes how to use those calculations for multiple writes.

Topics include:

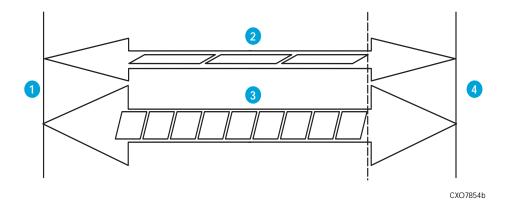
- Maximum number of messages, page 23
- Peak write rate, page 24
- Effects of multiple writes, page 27
- Actual I/Os per second, page 28

Maximum number of messages

A high bandwidth link can accommodate more data than a low bandwidth link (Figure 8). The maximum capacity is called the *bandwidth-latency product*. To calculate the maximum number of simultaneous messages allowed in the communications link:

- Multiply the net performance of the communications link (the total bits per second, minus overhead) by the one-way intersite latency (in seconds).
- Convert the result into bytes (Fibre Channel uses 10 bits per byte) and divide by the average message size (in bytes).

In calculating this number, it is assumed that the application can issue I/O as soon as space is available on the link, thus streaming the data from the host to the source array. This data is then replicated to the destination EVA with acknowledgment back to the host. It is also assumed that because there is only one source, data is written regularly and consistently. These assumptions are necessary to understand peak performance.



- 1 Site A
- 2 T3 link
- 3 OC3 link
- 4 Site B

Figure 8 Multiple writes for a single application

Peak write rate

In Figure 8, the parallelograms represent multiple write data streams from the same application. Each data stream contains data packets of equal size (same number of bytes). The low bandwidth link can accommodate three packets, but the high bandwidth link can accommodate nine packets. It is assumed that each new packet immediately follows the preceding packet, regardless of link bandwidth.

To determine the approximate peak write rate for an application:

- 1. Invert the single write percentage of bandwidth used (Table 5 on page 15) and divide by 100. The result is the maximum number of data streams that can be supported by that link for that distance.
- 2. Multiply the result by the single write rate.

Limitations

Between any two array controllers in a replication relationship, a bidirectional tunnel is created to transport the data (full copies and outstanding writes). A tunnel is a path that exists within the intersite link between two controllers and supports independent data transfers. With two controllers in each array, the four controllers create four tunnels (two active and two standby). A DR group belongs to only one tunnel, but one tunnel can support multiple DR groups.

The number of replication relationships required determines the total number of tunnels that are created (four for each relationship). Further, the number of tunnels determines the minimum required bandwidth for a specific communications link. For example, you want to consolidate two sources into one destination supporting eight tunnels (four active and four standby). Add the throughput requirements for each active tunnel using a given intersite link technology. This sum (the total throughput requirements for all tunnels) determines the maximum required bandwidth for that communications link.

The following factors can limit the expected maximum number of writes in the link:

Raw capability of the host, FCA, and the source array controller

For example, to ensure write order across all members of a DR group, the members are restricted to using the same FCA port and array controller and Secure Path's load balancing is turned off.

Number of buffer-to-buffer credits allowed between two Fibre Channel devices

For each replication relationship, buffers process the replication I/Os. The number of buffers depends on the intersite latency. For example, if the intersite latency is small, not all buffers are used. As the intersite latency increases, the number of buffers used will increase, which may become a performance limitation.

Tables 6, 7, and 8 list the buffer-to-buffer credits for B-series, C-series, and M-series switches, respectively. This limit becomes the bottleneck on a long-distance direct Fibre Channel connection with very long-distance gigabit interface converters (GBICs) or a wavelength division multiplexing (WDM) solution. It is not usually seen in Continuous Access EVA over IP configurations because the credit is returned by the IP gateway to the sending switch.

Maximum number of outstanding writes the array controller can support

Using Virtual Controller Software version 3.00 or 3.01, the tunnel within the array controller allows up to 31 (8 KB) outstanding writes per controller port. If the bandwidth latency product of the link can support more than 31 writes (such as in high speed, very long distance Continuous Access EVA over IP configurations), the maximum number of 8 KB or smaller outstanding writes in the link will be 31 for each array controller sharing the link.

Using Virtual Controller Software version 3.02 or later, up to 62 outstanding writes are allowed for each controller port. In a worst-case situation in which one port of each controller uses the same link, the combined limit is 124 (62 for each port) outstanding writes, subject to other limits.

Impact of the initial full copy process on the intersite link

In all 3.x versions of Virtual Controller Software, the full copy process allocates eight 128 KB data buffers. These buffers are only used during the initial full copy (a sequential read/write process in which a complete copy of the source array's content is made on the destination array) or after the write history log becomes full. When the log is full and the DR group members are marked for full copy, either a full copy or a fast resynchronization occurs. Fast resynchronization is the process of only moving those blocks of data that changed instead of the entire contents of the DR group members. This process uses random reads to find the data that must be copied to the destination array.

Replication consumes as much as 492 KB (62×8 KB), but the full copy consumes as much as 1 MB (8×128 KB). Combined, these two processes may overwhelm an intersite link if it does not have enough bandwidth.

Table 6 FC buffer-to-buffer credits for B-series switches

Switch family	Default credits	Credits available with Extended Fabric License
3xxx at 1 Gb/s	16	60
4xxx at 1 Gb/s	27	60
3xxx at 2 Gb/s	27	64
4xxx at 2 Gb/s	27	64

 Table 7
 FC buffer-to-buffer credits for C-series switches

Switch family	Default credits	Maximum available credits
All 16-port modules, Fx mode	16	255
E or TE modes	255	255
All 32-port modules	12	N/A

 Table 8
 FC buffer-to-buffer credits for M-series switches

Switch family	Default credits	Maximum available credits
Зххх	60	60
4xxx	162	162
бххх	60	60

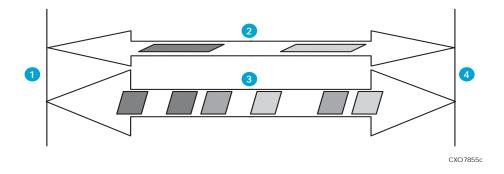
Effects of multiple writes

Figure 9 represents multiple writes from multiple applications. The difference between Figure 9 and the single application in Figure 8 is that the space between the writes is wider for multiple applications. The wider space reduces the number of writes in the intersite link and also reduces the maximum utilization rate of that link. Any differences in the number of writes in the link are based on how the bandwidth is shared between the applications.



TIP: If you are using multiple high-performance applications, HP recommends that you add Continuous Access EVA storage arrays to maintain performance.

In real-world application environments, it is not possible for all writes to arrive at precisely the expected times. Independent applications do not coordinate with each other when sending reads or writes to the source array. This lack of coordination creates space between any two writes on the intersite link. Mathematical queue theory suggests that the expected peak utilization is approximately 70% and the expected average is approximately 50%. In addition, there must be sufficient bandwidth on both fabrics for all traffic, if one fabric fails. If there is an even split between the two fabrics, the 50% rate becomes 25%.



- 1 Site A
- 2 T3 link
- 3 OC3 link
- 4 Site B

Figure 9 Multiple writes for multiple applications

To determine the effect of multiple writes on the intersite link, consider the following example. There are two applications—one performs 32 KB writes and one performs 2 KB writes. Table 5 on page 15 shows the bandwidth used for a 32 KB write. Table 9, using the same formula as in Table 5, shows the bandwidth used for a 2 KB write.

Table 9 Bandwidth used for a 2 KB write

Intersite link technology	Approximate IOPS	Throughput (Mb/s)	Approximate single stream bandwidth used
2 Gb/s Fibre Channel	417.50	8.35	0.4%
1 Gb/s Fibre Channel	405.60	8.11	0.8%
1 GbE (Gigabit Ethernet) IP	403.13	8.06	0.8%
OC3 IP	393.44	7.87	5.1%
E4 IP	391.97	7.84	5.6%
100 Mb/s IP	386.01	7.72	7.7%
T3 IP	361.40	7.23	16.1%
E3 IP	348.77	6.98	20.5%
10 Mb/s IP	246.08	4.92	49.2%
E1 IP	82.91	1.66	82.9%
T1 IP	64.77	1.30	86.4%

Actual I/Os per second

Determine the actual number of IOPS expected from each application by scaling the single write rate (either up or down) to the expected rate, and then scaling the bandwidth needed by the same number. If the total required bandwidth exceeds 25%, you cannot expect that link to support the replication requirements. For example, one application read and write produces transactions at 3 times the 32 KB single stream rates. The other read and write produces transactions at 4 times the 2 KB single stream rates.

Table 10 shows that only the 2 Gb/s link can support the replication requirements. If high-speed links are not available, you can use multiple low-speed links to obtain the required 25% of available bandwidth.

Table 10 Total bandwidth required for multiple I/O streams

Intersite link technology	32 KB IOPS (x3)	32 KB throughput (Mb/s)	Bandwidth used	2 KB IOPS (x4)	2 KB throughput (Mb/s)	Bandwidth used	Total bandwidth required
2 Gb/s Fibre Channel	938	300	15%	1670	33	2%	17%
1 Gb/s Fibre Channel	867	277	28%	1623	32	3%	31%
1 GbE (Gigabit Ethernet) IP	859	274	27%	1613	32	3%	31%
OC3 IP	623	199	129%	1574	31	20%	149%
E4 IP	599	192	138%	1568	31	23%	161%
100 Mb/s IP	522	167	167%	1544	31	31%	198%
T3 IP	334	107	238%	1446	29	64%	302.%
E3 IP	276	88	260%	1395	28	82%	342%
10 Mb/s IP	98	31	313%	984	20	197%	510%
E1 IP	21	6.6	330%	332	6.6	332%	662%
T1 IP	15	5	330%	259	5.2	345%	676%

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CAUTION: Design a solution that will not overload a controller or intersite link during normal operations to prevent significant loss of performance during failure conditions. This is especially true when using asynchronous replication and during the initial full copy or normalization process that occurs when you create a new DR group.

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